



Digital Lines & Continued Fractions

UPPSALA
UNIVERSITET

Hanna Uscka-Wehlou (Sweden)

A new recursive, continued fraction based, formula for digital lines $y = ax$ with irrational slopes $a = [0 ; a_1, a_2, a_3, \dots]$

The **index jump function** corresponding to a ; **essential 1's** :

$$\begin{array}{cccccccccccc}
 & b_1 & b_2 & \overbrace{1, 1}^{b_3} & b_4 & \overbrace{1, 1}^{b_5} & b_6 & b_7 & \overbrace{1, a_{11}}^{b_8} & b_9 & \overbrace{1, 1}^{b_{10}} & \overbrace{1, a_{16}}^{b_{11}} & b_{12} & \dots \\
 [0; \underline{1}, a_2, \underline{1}, 1, a_5, \underline{1}, 1, a_8, a_9, \underline{1}, a_{11}, a_{12}, \underline{1}, 1, \underline{1}, a_{16}, a_{17}, \dots] \\
 \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \dots \\
 (1, 2, 3, 5, 6, 8, 9, 10, 12, 13, 15, 17, \dots)
 \end{array}$$

The role of **essential 1's** in the construction of runs :

- $a_{i_a}(k+1) \geq 2 \Rightarrow S_k$ is the main run _{k} ,
- $a_{i_a}(k+1) = 1 \Rightarrow L_k$ is the main run _{k} .



The pixel-wise length of digital line segments

The **denominators** of the **convergents** :

For the **convergents** $\frac{p_n}{q_n} = [0; a_1, a_2, \dots, a_n]$ we have:
 $q_0 = 1$, $q_1 = a_1$, and, for $n \geq 2$, $q_n = a_n q_{n-1} + q_{n-2}$.

The **pixel-wise length** of runs on level n for the line $y = ax$:

$$|S_n| = q_{i_a(n+1)-1}, \quad |L_n| = q_{i_a(n+1)-1} + q_{i_a(n+1)-2}$$



Digital Lines & Continued Fractions

UPPSALA
UNIVERSITET

Hanna Uscka-Wehlou (Sweden)

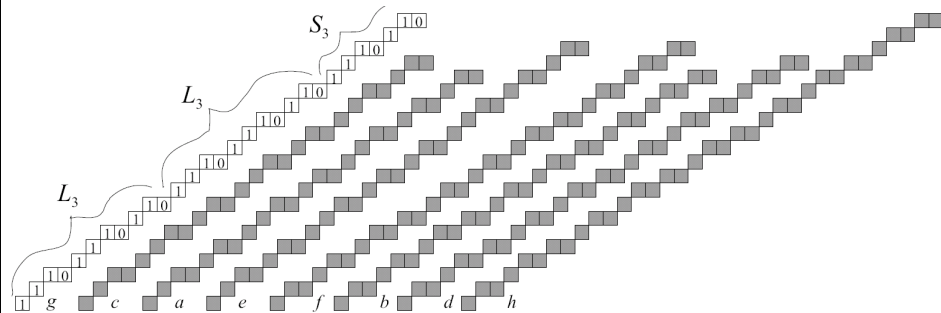
Two equivalence relations on the set of slopes

★ Quantitative (length)

★ Defined by run lengths (their cardinality)

All lines from the same class have the same run lengths on all digitization levels.

An example (from paper 3) :



$$\|S_1\|=1, \|S_2\|=2, \|S_3\|=2, \|S_4\|=3.$$

★ Qualitative (construction)

★ Defined by the places of essential 1's

All lines from the same class have the same construction in terms of long and short runs on all digitization levels.

Theorem (from paper 3) :

$$\forall n \in \mathbf{N}^+ \left[(\forall k \in [1, n-1]_{\mathbf{Z}} \ s_k = 2k) \right. \\ \left. \wedge (s_n > 2n \vee |J| = n-1) \right] \\ \Rightarrow \sup\{a; a \in [(s_j)_{j \in J}]_{\sim_{\text{con}}}\} = \frac{F_{2n-1}}{F_{2n}},$$

where $(F_n)_{n \in \mathbf{N}^+}$ is the **Fibonacci** sequence and $(s_j)_{j \in J}$ is the sequence of the places of essential 1's.